
Temporal Variability in Implicit Online Learning

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Abstract

In the setting of online learning, Implicit algorithms turn out to be highly successful from a practical standpoint. However, the tightest regret analyses only show marginal improvements over Online Mirror Descent. In this work, we shed light on this behavior carrying out a careful regret analysis. We prove a novel static regret bound that depends on the temporal variability of the sequence of loss functions, a quantity which is often encountered when considering dynamic competitors. We show, for example, that the regret can be constant if the temporal variability is constant and the learning rate is tuned appropriately, without the need of smooth losses. Moreover, we present an adaptive algorithm that achieves this regret bound without prior knowledge of the temporal variability and prove a matching lower bound. Finally, we validate our theoretical findings on classification and regression datasets.

1 Introduction

The *online learning* paradigm is a powerful tool to model common scenarios in the real world when the data comes in a streaming fashion, for example in the case of time series. In the last two decades there has been a tremendous amount of progress in this field (see, e.g., [30, 13, 24], for an introduction), which also led to advances in seemingly unrelated areas of machine learning and computer science. In this setting, a learning agent faces the environment in a game played sequentially. The protocol is the following: given a time horizon T , in every round $t = 1, \dots, T$ the agent chooses a model \mathbf{x}_t from a convex set V . Then, a convex loss function ℓ_t is revealed by the environment and the agent pays a loss $\ell_t(\mathbf{x}_t)$. As usual in this setting, we do not make assumptions about the environment, but allow it to be adversarial. The agent’s goal is to minimize her regret against any decision maker, i.e., the cumulative sum of her losses compared to the losses of an agent which always commits to the same choice \mathbf{u} . So, formally the regret against any $\mathbf{u} \in V$ is defined as

$$R_T(\mathbf{u}) \triangleq \sum_{t=1}^T \ell_t(\mathbf{x}_t) - \sum_{t=1}^T \ell_t(\mathbf{u}).$$

Much of the progress in this field is driven by the strictly related model of Online Linear Optimization (OLO): exploiting the assumption that the loss functions are convex, we can linearize them using a first-order approximation through its (sub)gradient and subsequently minimize the linearized regret. For example, the well-known Online Gradient Descent (OGD) [38] simply uses the direction of the negative (sub)gradient of the loss function to update its model, multiplied by a given learning rate. Usually, a properly tuned learning rate gives a regret bound of $\mathcal{O}(\sqrt{T})$, which is also optimal. On the other hand, we can choose to not use any approximation to the loss function and instead update our model using directly the loss function rather than its subgradient [17]. This type of update is known as *Implicit* and algorithms designed in this way are known to have practical advantages [18]. Unfortunately, their theoretical understanding is still limited at this point.

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Our first contribution (Section 5) in this paper is a refined analysis of Implicit algorithms in the framework of Online Mirror Descent (OMD). Doing this allows us to understand why Implicit algorithms might practically work better compared to algorithms which use (sub)gradients in the update. In particular, we describe how these algorithms can potentially incur only a constant regret if the sequence of loss functions does not vary with time. In particular, we measure the hardness of the sequence of loss functions with its *temporal variability*, which is defined as

$$V_T \triangleq \sum_{t=2}^T \max_{\mathbf{x} \in V} \ell_t(\mathbf{x}) - \ell_{t-1}(\mathbf{x}). \quad (1)$$

Our second contribution (Section 6) is a new adaptive Implicit algorithm, AdaImplicit, which retains the worst-case $\mathcal{O}(\sqrt{T})$ regret bound but takes advantage of a slow varying sequence of loss functions and achieve a regret of $\mathcal{O}(V_T + 1)$. Also, we prove a lower bound which shows that our algorithm is optimal. Finally, in order to show the benefits of using Implicit algorithms in practice, in Section 7 we conduct an empirical analysis on real-world datasets in both classification and regression tasks.

2 Related Work

Implicit Updates. The implicit updates in online learning were proposed for the first time by Kivinen and Warmuth [17]. However, such update with the Euclidean divergence is the Proximal update in the optimization literature dating back at least to 1965 [22, 19, 29, 27], and more recently used even in the stochastic setting [33, 2]. Later, this idea was re-invented by Crammer et al. [11] for the specific case of linear prediction with losses that have a range of values in which they are zero, e.g., hinge loss and epsilon-insensitive loss. Implicit updates were also used for online learning with kernels [9] and to deal with importance weights [16]. Kulis and Bartlett [18] provide the first regret bounds for implicit updates that match those of OMD, while McMahan [20] makes the first attempt to quantify the advantage of the implicit updates in the regret bound. Finally, Song et al. [31] generalize the results in McMahan [20] to Bregman divergences and strongly convex functions, and quantify the gain differently in the regret bound. Note that in [20, 31] the gain cannot be exactly quantified, providing just a non-negative data-dependent quantity subtracted to the regret bound.

Adaptivity. Our new analysis hinges on the concept of *temporal variability* V_T of the losses, a quantity first defined in Besbes et al. [5] in the context of non-stationary stochastic optimization and later generalized in Chen et al. [8]. In general, the temporal variability has been used in works considering dynamic environments [e.g., 15, 37, 3, 36]. In particular, Jadbabaie et al. [15] consider different notions of adaptivity at the same time: if we consider the static regret case with no optimistic updates, then their bound gives $R_T = \tilde{\mathcal{O}}(\sqrt{\sum_{t=1}^T \|g_t\|_*^2} + 1)$, which is never better than ours. At first sight, our algorithm seems to achieve the same constant regret bound of Optimistic algorithms [10, 28] if the sequence of loss functions is such that $V_T = \mathcal{O}(1)$. However, for this result Optimistic algorithms need either smooth or linear loss functions. In contrast, our algorithm does not need this assumption. Other examples of adaptivity to the sequence of loss functions can be found in [14, 32], which consider bounds in terms of the variance of the sequence of linear losses.

Finally, it is worth mentioning that recently there have been attempts to analyze Implicit algorithms in dynamic environments [see, e.g., 12, 1, 7]. Nevertheless, these works are not directly comparable to ours since they either consider a different (noisy) setting and competitor or make stronger assumptions (i.e. smoothness and/or strong convexity of the loss functions).

3 Definitions

For a function $f : \mathbb{R}^d \rightarrow (-\infty, +\infty]$, we define a *subgradient* of f in $\mathbf{x} \in \mathbb{R}^d$ as a vector $\mathbf{g} \in \mathbb{R}^d$ that satisfies $f(\mathbf{y}) \geq f(\mathbf{x}) + \langle \mathbf{g}, \mathbf{y} - \mathbf{x} \rangle$, $\forall \mathbf{y} \in \mathbb{R}^d$. We denote the set of subgradients of f in \mathbf{x} by $\partial f(\mathbf{x})$. The *indicator function of the set* V , $i_V : \mathbb{R}^d \rightarrow (-\infty, +\infty]$, is defined as

$$i_V(\mathbf{x}) = \begin{cases} 0, & \mathbf{x} \in V, \\ +\infty, & \text{otherwise.} \end{cases}$$

We denote the *dual norm* of $\|\cdot\|$ by $\|\cdot\|_*$. A proper function $f : \mathbb{R}^d \rightarrow (-\infty, +\infty]$ is μ -strongly convex over a convex set $V \subseteq \text{int dom } f$ w.r.t. $\|\cdot\|$ if $\forall \mathbf{x}, \mathbf{y} \in V$ and $\mathbf{g} \in \partial f(\mathbf{x})$, we have

Algorithm 1 Implicit Online Mirror Descent (IOMD)

Require: Non-empty closed convex set $V \subset X \subset \mathbb{R}^d$, $\psi : X \rightarrow \mathbb{R}$, $\eta_t > 0$, $\mathbf{x}_1 \in V$

- 1: **for** $t = 1, \dots, T$ **do**
 - 2: Output $\mathbf{x}_t \in V$
 - 3: Receive $\ell_t : \mathbb{R}^d \rightarrow \mathbb{R}$ and pay $\ell_t(\mathbf{x}_t)$
 - 4: Update $\mathbf{x}_{t+1} = \arg \min_{\mathbf{x} \in V} B_\psi(\mathbf{x}, \mathbf{x}_t) + \eta_t \ell_t(\mathbf{x})$
 - 5: **end for**
-

$f(\mathbf{y}) \geq f(\mathbf{x}) + \langle \mathbf{g}, \mathbf{y} - \mathbf{x} \rangle + \frac{\mu}{2} \|\mathbf{x} - \mathbf{y}\|^2$. Let $\psi : X \rightarrow \mathbb{R}$ be strictly convex and continuously differentiable on $\text{int } X$. The *Bregman Divergence* w.r.t. ψ is $B_\psi : X \times \text{int } X \rightarrow \mathbb{R}_+$ defined as $B_\psi(\mathbf{x}, \mathbf{y}) = \psi(\mathbf{x}) - \psi(\mathbf{y}) - \langle \nabla \psi(\mathbf{y}), \mathbf{x} - \mathbf{y} \rangle$. We assume that ψ is strongly convex w.r.t. a norm $\|\cdot\|$ in $\text{int } X$. We also assume w.l.o.g. the strong convexity constant to be 1, which implies

$$B_\psi(\mathbf{x}, \mathbf{y}) \geq \frac{1}{2} \|\mathbf{x} - \mathbf{y}\|^2, \quad \forall \mathbf{x} \in X, \mathbf{y} \in \text{int } X. \quad (2)$$

4 Online Mirror Descent with Implicit Updates

In this section, we introduce the Implicit Online Mirror Descent (IOMD) algorithm, its relationship with OMD, and some of its properties.

Consider a set $V \subset X \subseteq \mathbb{R}^d$. The Online Mirror Descent [35, 4] update over V is

$$\mathbf{x}_{t+1} = \arg \min_{\mathbf{x} \in V} B_\psi(\mathbf{x}, \mathbf{x}_t) + \eta_t (\ell_t(\mathbf{x}_t) + \langle \mathbf{g}_t, \mathbf{x} - \mathbf{x}_t \rangle) = \arg \min_{\mathbf{x} \in V} B_\psi(\mathbf{x}, \mathbf{x}_t) + \eta_t \langle \mathbf{g}_t, \mathbf{x} \rangle,$$

for $\mathbf{g}_t \in \partial \ell_t(\mathbf{x}_t)$ received as feedback. In words, OMD updates the solution minimizing a first-order approximation of the received loss, ℓ_t , around the predicted point, \mathbf{x}_t , constrained to be not too far from the predicted point measured with the Bregman divergence. It is well-known, [e.g. 24], that the regret guarantee for OMD for a non-increasing sequence of learning rates $(\eta_t)_{t=1}^T$ is

$$R_T(\mathbf{u}) \leq \sum_{t=1}^T \frac{B_\psi(\mathbf{u}, \mathbf{x}_t) - B_\psi(\mathbf{u}, \mathbf{x}_{t+1})}{\eta_t} + \sum_{t=1}^T \frac{\eta_t}{2} \|\mathbf{g}_t\|_*^2, \quad \forall \mathbf{u} \in V. \quad (3)$$

This gives a $\mathcal{O}(\sqrt{T})$ regret with, e.g., $\max_{\mathbf{x}, \mathbf{y} \in V} B_\psi(\mathbf{x}, \mathbf{y}) < \infty$, Lipschitz losses, and $\eta_t \propto 1/\sqrt{t}$.

A natural variation of the classic OMD update is to use the actual loss function ℓ_t , rather than its first-order approximation. This is called *implicit update* [17] and is defined as

$$\mathbf{x}_{t+1} = \arg \min_{\mathbf{x} \in V} B_\psi(\mathbf{x}, \mathbf{x}_t) + \eta_t \ell_t(\mathbf{x}). \quad (4)$$

Note that, in general, this update does not have a closed form, but for many interesting cases it is still possible to efficiently compute it. Notably, for $\psi = \frac{1}{2} \|\cdot\|_2^2$ and linear prediction with the square, absolute, and hinge loss, these updates can all be computed in closed form when $V = \mathbb{R}^d$ [see, e.g., 11, 18]. This update leads to the Implicit Online Mirror Descent (IOMD) algorithm in Algorithm 1.

We next show how the update in Eq. (4) yields new interesting properties which are not shared with its non-implicit counterpart. Their proofs can be found in Appendix B.

Proposition 4.1. *Let \mathbf{x}_{t+1} be defined as in Eq. (4). Then, there exists $\mathbf{g}'_t \in \partial \ell_t(\mathbf{x}_{t+1})$ such that*

$$\ell_t(\mathbf{x}_t) - \ell_t(\mathbf{x}_{t+1}) - B_\psi(\mathbf{x}_{t+1}, \mathbf{x}_t)/\eta_t \geq 0, \quad (5)$$

$$\langle \eta_t \mathbf{g}'_t + \nabla \psi(\mathbf{x}_{t+1}) - \nabla \psi(\mathbf{x}_t), \mathbf{u} - \mathbf{x}_{t+1} \rangle \geq 0 \quad \forall \mathbf{u} \in V, \quad (6)$$

$$\langle \mathbf{g}'_t, \mathbf{x}_{t+1} - \mathbf{x}_t \rangle \geq \langle \mathbf{g}_t, \mathbf{x}_{t+1} - \mathbf{x}_t \rangle. \quad (7)$$

The first property implies that, in contrast to OMD, the value of the loss function in \mathbf{x}_{t+1} is always smaller than or equal to its value in \mathbf{x}_t . This means that, if $\ell_t = \ell$, the value $\ell(\mathbf{x}_t)$ will be monotonically decreasing over time. The second property gives an alternative way to write the update rule expressed in Eq. (4). In particular, using $\psi(\mathbf{x}) = \frac{1}{2} \|\mathbf{x}\|_2^2$ and $V = \mathbb{R}^d$ the update becomes $\mathbf{x}_{t+1} = \mathbf{x}_t - \eta_t \mathbf{g}'_t$, motivating the name “implicit”. Using this fact in the last property,² we have that with L_2 regularization, the dual norm of \mathbf{g}'_t is smaller than the dual norm of \mathbf{g}_t , i.e. $\|\mathbf{g}'_t\|_2 \leq \|\mathbf{g}_t\|_2$.

²Eq. (7) is nothing else than the fact that subgradients are monotone operators.

Let's gain some additional intuition on the implicit updates. Consider the case of $V = \mathbb{R}^d$ and $\psi(\mathbf{x}) = \frac{1}{2} \|\cdot\|_2^2$. We have that $\mathbf{x}_{t+1} = \mathbf{x}_t - \eta_t \mathbf{g}'_t$, where $\mathbf{g}'_t \in \partial \ell_t(\mathbf{x}_{t+1})$. Now, if $\ell_{t+1} \approx \ell_t$, we would be updating the algorithm approximately with the *next subgradient*. On the other hand, knowing future gradients is a safe way to have constant regret. Hence, we can expect IOMD to have low regret if the functions are slowly varying over time. In the next sections, we will see that this is indeed the case.

5 Two Regret Bounds for IOMD

In the following, we will present a new regret guarantee for IOMD. First, we give a simple lemma that provides a bound on the cumulative losses paid *after* the updates (proof in Appendix B).

Lemma 5.1. *Let $V \subset X \subset \mathbb{R}^d$ be a non-empty closed convex set. Let B_ψ be the Bregman divergence w.r.t. $\psi : X \rightarrow \mathbb{R}$. Then, Algorithm 1 guarantees*

$$\sum_{t=1}^T \ell_t(\mathbf{x}_{t+1}) - \sum_{t=1}^T \ell_t(\mathbf{u}) \leq \sum_{t=1}^T \frac{B_\psi(\mathbf{u}, \mathbf{x}_t) - B_\psi(\mathbf{u}, \mathbf{x}_{t+1})}{\eta_t} - \sum_{t=1}^T \frac{1}{\eta_t} B_\psi(\mathbf{x}_{t+1}, \mathbf{x}_t). \quad (8)$$

Furthermore, assume that $(\eta_t)_{t=1}^T$ is a non-increasing sequence and let $D^2 \triangleq \max_{\mathbf{x}, \mathbf{u} \in V} B_\psi(\mathbf{u}, \mathbf{x})$. Then the bound can further be expressed as

$$\sum_{t=1}^T \ell_t(\mathbf{x}_{t+1}) - \sum_{t=1}^T \ell_t(\mathbf{u}) \leq \frac{D^2}{\eta_T} - \sum_{t=1}^T \frac{1}{\eta_t} B_\psi(\mathbf{x}_{t+1}, \mathbf{x}_t). \quad (9)$$

Adding $\sum_{t=1}^T \ell_t(\mathbf{x}_t)$ on both sides of Eq. (8), we immediately get our new regret bound.

Theorem 5.2. *Under the assumptions of Lemma 5.1, the regret incurred by Algorithm 1 is bounded as*

$$R_T(\mathbf{u}) \leq \sum_{t=1}^T \frac{B_\psi(\mathbf{u}, \mathbf{x}_t) - B_\psi(\mathbf{u}, \mathbf{x}_{t+1})}{\eta_t} + \sum_{t=1}^T \left[\ell_t(\mathbf{x}_t) - \ell_t(\mathbf{x}_{t+1}) - \frac{B_\psi(\mathbf{x}_{t+1}, \mathbf{x}_t)}{\eta_t} \right]. \quad (10)$$

We note that this result could also be extrapolated from [31], by carefully going through the proof of their Lemma 1. However, as in the other previous work, they did not identify that the key quantity to be used in order to quantify an actual gain is the temporal variability V_T , as we will show later.

First Regret: Recovering OMD's Guarantee. To this point, the advantages of an implicit update are still not clear. Therefore, we now show how, from Theorem 5.2, one can get a possibly tighter bound than the usual $\mathcal{O}(\sqrt{T})$. The key point in this new analysis is to introduce \mathbf{g}'_t as defined in Proposition 4.1 and relate it to the Bregman divergence between \mathbf{x}_t and \mathbf{x}_{t+1} .

Theorem 5.3. *Let $\mathbf{g}'_t \in \partial \ell_t(\mathbf{x}_{t+1})$ satisfy Eq. (6). Assume ψ to be 1-strongly convex w.r.t. $\|\cdot\|$. Then, under the assumptions of Lemma 5.1, we have that Algorithm 1 satisfies*

$$\ell_t(\mathbf{x}_t) - \ell_t(\mathbf{x}_{t+1}) - \frac{B_\psi(\mathbf{x}_{t+1}, \mathbf{x}_t)}{\eta_t} \leq \eta_t \|\mathbf{g}_t\|_* \min \left(2 \|\mathbf{g}'_t\|_*, \frac{\|\mathbf{g}_t\|_*}{2} \right), \quad \forall t, \mathbf{g}_t \in \partial \ell_t(\mathbf{x}_t). \quad (11)$$

Proof. Using the convexity of the losses, we can bound the difference between $\ell_t(\mathbf{x}_t)$ and $\ell_t(\mathbf{x}_{t+1})$:

$$\ell_t(\mathbf{x}_t) - \ell_t(\mathbf{x}_{t+1}) \leq \langle \mathbf{g}_t, \mathbf{x}_t - \mathbf{x}_{t+1} \rangle \leq \|\mathbf{g}_t\|_* \|\mathbf{x}_t - \mathbf{x}_{t+1}\|,$$

where $\mathbf{g}_t \in \partial \ell_t(\mathbf{x}_t)$. Given that ψ is 1-strongly convex, we can use Eq. (2) to obtain

$$\ell_t(\mathbf{x}_t) - \ell_t(\mathbf{x}_{t+1}) \leq \|\mathbf{g}_t\|_* \sqrt{2B_\psi(\mathbf{x}_{t+1}, \mathbf{x}_t)}. \quad (12)$$

Note that $\ell_t(\mathbf{x}_t) - \ell_t(\mathbf{x}_{t+1}) - B_\psi(\mathbf{x}_{t+1}, \mathbf{x}_t)/\eta_t \leq \ell_t(\mathbf{x}_t) - \ell_t(\mathbf{x}_{t+1})$. Hence, to get the first term in the min of Eq. (11), we can simply look for an upper bound on the term $\sqrt{2B_\psi(\mathbf{x}_{t+1}, \mathbf{x}_t)}$ in Eq. (12) above. Using the fact that the Bregman divergence is convex in its first argument, we get

$$\begin{aligned} B_\psi(\mathbf{x}_{t+1}, \mathbf{x}_t) &\leq \langle \nabla \psi(\mathbf{x}_{t+1}) - \nabla \psi(\mathbf{x}_t), \mathbf{x}_{t+1} - \mathbf{x}_t \rangle \leq \langle \eta_t \mathbf{g}'_t, \mathbf{x}_t - \mathbf{x}_{t+1} \rangle \leq \eta_t \|\mathbf{g}'_t\|_* \|\mathbf{x}_{t+1} - \mathbf{x}_t\| \\ &\leq \eta_t \|\mathbf{g}'_t\|_* \sqrt{2B_\psi(\mathbf{x}_{t+1}, \mathbf{x}_t)}, \end{aligned}$$

where we used Eq. (6) in the second inequality and Eq. (2) in the last one. Solving this inequality with respect to $B_\psi(\mathbf{x}_{t+1}, \mathbf{x}_t)$, we get $\sqrt{2B_\psi(\mathbf{x}_{t+1}, \mathbf{x}_t)} \leq 2\eta_t \|\mathbf{g}'_t\|_*$.

For the second term, it suffices to subtract $B_\psi(\mathbf{x}_{t+1}, \mathbf{x}_t)/\eta_t$ on both sides of Eq. (12) and use the fact that $bx - \frac{a}{2}x^2 \leq \frac{b^2}{2a}, \forall x \in \mathbb{R}$ with $x = B_\psi(\mathbf{x}_{t+1}, \mathbf{x}_t)$. \square

This Theorem immediately gives us that Algorithm 1 has a regret upper-bounded by

$$R_T(\mathbf{u}) \leq \sum_{t=1}^T \frac{B_\psi(\mathbf{u}, \mathbf{x}_t) - B_\psi(\mathbf{u}, \mathbf{x}_{t+1})}{\eta_t} + \sum_{t=1}^T \eta_t \|\mathbf{g}_t\|_* \min \left(2\|\mathbf{g}'_t\|_*, \frac{\|\mathbf{g}_t\|_*}{2} \right), \quad (13)$$

where $\mathbf{g}_t \in \partial \ell_t(\mathbf{x}_t)$. The presence of the minimum makes this bound equivalent in a worst-case sense to the one of OMD in Eq. (3). Moreover, at least in the Euclidean case, from Eq. (7) we have that $\|\mathbf{g}'_t\|_2 \leq \|\mathbf{g}_t\|_2$. However, it is difficult to quantify the gain over OMD because in general $\|\mathbf{g}_t\|_*$ and $\|\mathbf{g}'_t\|_*$ are data-dependent. Hence, as in the other previous analyses, the gain over OMD would be only marginal and not quantifiable. This is not a limit of our analysis: it is easy to realize that in the worst case the OMD update and the IOMD update can coincide. To show instead that a real gain is possible, we are now going to take a different path.

Second Regret: Temporal Variability in IOMD. Here we formalize our key intuition that IOMD is using an approximation of the future subgradient when the losses do not vary much over time. We use the notion of *temporal variability of the losses*, V_T , as given in Eq. (1). Considering again our regret bound in Theorem 5.2 and using $\eta_t = \eta$ for all t , we immediately have

$$\begin{aligned} R_T(\mathbf{u}) &\leq \frac{B_\psi(\mathbf{u}, \mathbf{x}_1)}{\eta} + \sum_{t=1}^T \left(\ell_t(\mathbf{x}_t) - \ell_t(\mathbf{x}_{t+1}) - \frac{B_\psi(\mathbf{x}_{t+1}, \mathbf{x}_t)}{\eta} \right) \\ &\leq \frac{B_\psi(\mathbf{u}, \mathbf{x}_1)}{\eta} + \ell_1(\mathbf{x}_1) - \ell_T(\mathbf{x}_{T+1}) + \sum_{t=2}^T \left(\max_{\mathbf{x} \in V} \ell_t(\mathbf{x}) - \ell_{t-1}(\mathbf{x}) - \frac{B_\psi(\mathbf{x}_{t+1}, \mathbf{x}_t)}{\eta} \right) \\ &\leq \frac{B_\psi(\mathbf{u}, \mathbf{x}_1)}{\eta} + \ell_1(\mathbf{x}_1) - \ell_T(\mathbf{x}_{T+1}) + V_T. \end{aligned}$$

This means that using a *constant learning rate yields a regret bound of $\mathcal{O}(V_T + 1)$, which might be better than $\mathcal{O}(\sqrt{T})$ if the temporal variability is low*. In particular, we can even get constant regret if $V_T = \mathcal{O}(1)$. On the contrary, OMD cannot achieve a constant regret for any convex loss even if $V_T = 0$, since it would imply an impossible $\mathcal{O}(1/T)$ rate for non-smooth batch black-box optimization [23, Theorem 3.2.1]. Instead, IOMD does not violate the lower bound since it is not a black-box method. As far as we know, the connection between IOMD and temporal variability has never been observed before. On the other hand, even when the temporal variability is high, we can still use a $\mathcal{O}(1/\sqrt{T})$ learning rate to achieve a worst case regret of the order $\mathcal{O}(\sqrt{T})$.

We would like to point out that a similar behaviour arises from *Follow The Regularized Leader* algorithm (FTRL) employed with full losses, rather than linearized ones. We show a detailed derivation in Appendix E. Unfortunately, contrarily to the OMD case employing FTRL would entail solving a constrained convex optimization problem whose size (in terms of number of functions) grows each step, that would have a high running time even when the implicit updates have closed form expressions, e.g., linear classification with hinge loss.

Finally, a natural question arises: can we get a bound which interpolates between $\mathcal{O}(V_T + 1)$ and $\mathcal{O}(\sqrt{T})$, without any prior knowledge on the quantity V_T ? We give a positive answer to this question by presenting an adaptive strategy in the next section.

6 Adapting to the Temporal variability with AdaImplicit

In this section, we present an adaptive strategy to set the learning rates, in order to give a regret guarantee that depends optimally on the temporal variability.

From the previous section, we saw that the key quantity in the IOMD regret bound is

$$\delta_t \triangleq \ell_t(\mathbf{x}_t) - \ell_t(\mathbf{x}_{t+1}) - \frac{B_\psi(\mathbf{x}_{t+1}, \mathbf{x}_t)}{\eta_t}. \quad (14)$$

Algorithm 2 AdaImplicit

Require: Non-empty closed convex set $V \subset X \subset \mathbb{R}^d$, $\psi : X \rightarrow \mathbb{R}$, $\lambda_1 = 0$, $\beta^2 > 0$, $\mathbf{x}_1 \in V$

- 1: **for** $t = 1, \dots, T$ **do**
 - 2: Output $\mathbf{x}_t \in V$
 - 3: Receive $\ell_t : \mathbb{R}^d \rightarrow \mathbb{R}$ and pay $\ell_t(\mathbf{x}_t)$
 - 4: Update $\mathbf{x}_{t+1} = \arg \min_{\mathbf{x} \in V} \ell_t(\mathbf{x}) + \lambda_t B_\psi(\mathbf{x}, \mathbf{x}_t)$
 - 5: Set $\delta_t = \ell_t(\mathbf{x}_t) - \ell_t(\mathbf{x}_{t+1}) - \lambda_t B_\psi(\mathbf{x}_{t+1}, \mathbf{x}_t)$
 - 6: Update $\lambda_{t+1} = \lambda_t + \frac{1}{\beta^2} \delta_t$
 - 7: **end for**
-

From Eq. (5), we have that $\delta_t \geq 0$. At this point, one might think of using a doubling trick: monitor $\sum_{i=1}^t \delta_i$ over time and restart the algorithm with a different learning rate once it exceeds a certain threshold. In Appendix A, we show that it is indeed possible to use such a strategy. However, while theoretically effective, we can't expect the doubling trick to have any decent performance in practice. Consequently, we are going to show how to use instead an *adaptive* learning rate.

AdaImplicit. Define $D^2 \triangleq \max_{\mathbf{x}, \mathbf{u} \in V} B_\psi(\mathbf{u}, \mathbf{x})$ and assume $D < \infty$. For ease of notation, we let $\eta_t = 1/\lambda_t$ where λ_t will be decided in the following. Assuming $(\lambda_t)_{t=1}^T$ to be an increasing sequence, from Theorem 5.2 we get

$$R_T(\mathbf{u}) \leq D^2 \lambda_T + \sum_{t=1}^T [\ell_t(\mathbf{x}_t) - \ell_t(\mathbf{x}_{t+1}) - \lambda_t B_\psi(\mathbf{x}_{t+1}, \mathbf{x}_t)]. \quad (15)$$

Ideally, to minimize the regret we would like to have λ_T to be as close as possible to the sum over time in the r.h.s. of this expression. However, setting $\lambda_t \propto \sum_{s=1}^t \delta_i$ would introduce an annoying recurrence in the computation of λ_t . To solve this issue, we explore the same strategy adopted in AdaFTRL [25], adapting it to the OMD case: we set $\lambda_{t+1} = \frac{1}{\beta^2} \sum_{i=1}^t \delta_i$ for $t \geq 2$, for a parameter β to be defined later, and $\lambda_1 = 0$. We call the resulting algorithm AdaImplicit and describe it in Algorithm 2. Before proving a regret bound for it, we first provide a technical lemma for the analysis. This lemma can be found in [24, 26] and for completeness we give a proof in Appendix B.

Lemma 6.1. *Let $\{a_t\}_{t=1}^\infty$ be any sequence of non-negative real numbers. Suppose that $\{\Delta_t\}_{t=1}^\infty$ is a sequence of non-negative real numbers satisfying $\Delta_1 = 0$ and³ $\Delta_{t+1} \leq \Delta_t + \min\{ba_t, ca_t^2/(2\Delta_t)\}$, for any $t \geq 1$. Then, for any $T \geq 0$, $\Delta_{T+1} \leq \sqrt{(b^2 + c) \sum_{t=1}^T a_t^2}$.*

We are now ready to prove a regret bound for Algorithm 2.

Theorem 6.2. *Let $V \subset X \subset \mathbb{R}^d$ be a non-empty closed convex set. Let B_ψ be the Bregman divergence w.r.t. $\psi : X \rightarrow \mathbb{R}$ and let $D^2 = \max_{\mathbf{x}, \mathbf{u} \in V} B_\psi(\mathbf{u}, \mathbf{x})$. Assume ψ to be 1-strongly convex with respect to $\|\cdot\|$ in V . Then, for any $\mathbf{u} \in V$, running Algorithm 2 with $\beta = D$ guarantees*

$$R_T(\mathbf{u}) \leq \min \left\{ 2(\ell_1(\mathbf{x}_1) - \ell_T(\mathbf{x}_{T+1}) + V_T), 2D \sqrt{3 \sum_{t=1}^T \|\mathbf{g}_t\|_*^2} \right\}, \quad \forall \mathbf{g}_t \in \ell_t(\mathbf{x}_t). \quad (16)$$

Proof. Using the definition of λ_t and the fact that the sequence $(\lambda_t)_{t=1}^{T+1}$ is increasing over time, the regret in Eq. (15) can be upper bounded as $R_T(\mathbf{u}) \leq (D^2 + \beta^2) \lambda_{T+1}$. Therefore, we need an upper bound on λ_{T+1} . We split the proof in two parts, one for each term in the min in Eq. (16). For the first term, using the definition of λ_t we have

$$\begin{aligned} \beta^2 \lambda_{T+1} &= \sum_{t=1}^T [\ell_t(\mathbf{x}_t) - \ell_t(\mathbf{x}_{t+1}) - \lambda_t B_\psi(\mathbf{x}_{t+1}, \mathbf{x}_t)] \\ &\leq \ell_1(\mathbf{x}_1) - \ell_T(\mathbf{x}_{T+1}) + \sum_{t=2}^T [\ell_t(\mathbf{x}_t) - \ell_{t-1}(\mathbf{x}_t)] \leq \ell_1(\mathbf{x}_1) - \ell_T(\mathbf{x}_{T+1}) + V_T, \end{aligned}$$

³With a small abuse of notation, let $\min(x, y/0) = x$.

from which using $\beta = D$ the result follows.

For the second term, from Lemma 5.3 for $t \geq 2$ we have $\delta_t \leq \frac{\|\mathbf{g}_t\|_*^2}{2\lambda_t}$. On the other hand,

$$\begin{aligned} \delta_t &= \ell_t(\mathbf{x}_t) - \ell_t(\mathbf{x}_{t+1}) - \lambda_t B_\psi(\mathbf{x}_{t+1}, \mathbf{x}_t) \leq \ell_t(\mathbf{x}_t) - \ell_t(\mathbf{x}_{t+1}) \leq \langle \mathbf{g}_t, \mathbf{x}_t - \mathbf{x}_{t+1} \rangle \\ &\leq \|\mathbf{g}_t\|_* \|\mathbf{x}_t - \mathbf{x}_{t+1}\| \leq \sqrt{2D} \|\mathbf{g}_t\|_*, \end{aligned}$$

where in the last step we used Eq. (2) and the definition of D . Therefore, putting the last two results together we get

$$\delta_t \leq \min \left(\sqrt{2D} \|\mathbf{g}_t\|_*, \|\mathbf{g}_t\|_*^2 / (2\lambda_t) \right), \quad \forall \mathbf{g}_t \in \partial \ell_t(\mathbf{x}_t).$$

Note that $\lambda_{t+1} = \lambda_t + \frac{1}{\beta^2} \delta_t$. Hence, $\lambda_1 = 0$, $\lambda_2 = (\ell_1(\mathbf{x}_1) - \ell_1(\mathbf{x}_2)) / \beta^2 \leq \sqrt{2D} \|\mathbf{g}_1\|_* / \beta^2$, and

$$\lambda_{t+1} = \lambda_t + \frac{1}{\beta^2} \delta_t \leq \lambda_t + \frac{1}{\beta^2} \min \left(\sqrt{2D} \|\mathbf{g}_t\|_*, \frac{\|\mathbf{g}_t\|_*^2}{2\lambda_t} \right), \quad \forall t \geq 3.$$

Therefore, using Lemma 6.1 with $\Delta_t = \lambda_t$, $b = \frac{\sqrt{2D}}{\beta^2}$ and $c = \frac{1}{\beta^2}$, $a_t = \|\mathbf{g}_t\|_*$, we get

$$\lambda_{T+1} \leq \sqrt{(2D^2/\beta^4 + 1/\beta^2) \sum_{t=1}^T \|\mathbf{g}_t\|_*^2},$$

from which setting $\beta = D$ we obtain the second term in the min in Eq. (16). \square

This last theorem shows that Algorithm 2 can have a low regret if the temporal variability of the losses V_T is low. Moreover, differently from Optimistic Algorithms, Algorithm 2 does not need additional assumptions on the losses (for example smoothness), as done for example in [15].

Lower Bound. Next, we are going to prove a lower bound in terms of the temporal variability V_T , which shows that the regret bound in Theorem 6.2 cannot be improved further. The proof is a simple modification of the standard arguments used to prove lower bounds for constrained OLO and is reported in Appendix B.

Theorem 6.3. *Let $d \geq 2$, $\|\cdot\|$ an arbitrary norm on \mathbb{R}^d , and $V = \{\mathbf{x} \in \mathbb{R}^d : \|\mathbf{x}\| \leq D/2\}$. Let \mathcal{A} be a deterministic algorithm on V . Let T be any non-negative integer. Then, for any $V'_T \geq 0$, there exists a sequence of convex loss functions $\ell_1(\mathbf{x}), \dots, \ell_T(\mathbf{x})$ with temporal variability equal to V'_T and $\mathbf{u} \in V$ such that the regret of algorithm \mathcal{A} satisfies $R_T(\mathbf{u}) \geq V'_T$.*

7 Empirical results

In this section, we compare the empirical performance of our algorithm AdaImplicit with standard baselines in online learning:

OGD [38], OGD with adaptive learning rate $\eta_t = \frac{\beta}{\sqrt{\sum_{i=1}^t \|\mathbf{g}_i\|_*^2}}$ (AdaOGD) [21], and IOMD with $\eta_t = \beta/\sqrt{t}$ (Implicit) [18].

Synthetic Experiment. We first show the benefits of AdaImplicit on a synthetic dataset. The loss functions are chosen to have a small temporal variability V_T . In particular, we consider a 1- d case using $\ell_t(x) = \frac{1}{4}(x - y_t)^2$ with $y_t = 100 \sin(\pi \frac{t}{10T})$, a time horizon $T = 2000$ and the L_2 ball of diameter $D = 150$. We set $\beta = 1$ in all algorithms. The update of the implicit algorithms can be computed in closed form: $x_{t+1} = x_t - \frac{\eta_t}{2+\eta_t}(x_t - y_t)$. In Fig. 1 we show the cumulative loss $L_T = \sum_{t=1}^T \ell_t(x_t)$ of the algorithms (note that the y -axis is plotted in logarithmic scale). From the figure we can see that, contrarily to the other algorithms, the cumulative loss of AdaImplicit grows slowly over time, reflecting experimentally the bound given in Theorem 6.2. Also, even if not directly observable, OGD and IOMD basically incur the same total cumulative loss.

Real world datasets. We are now going to show some experiments conducted on real data. Here, there is no reason to believe that the temporal variability is small. However, we still want to verify

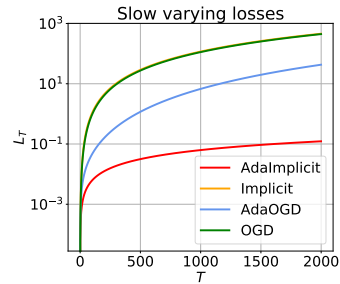


Figure 1: Synthetic experiment.

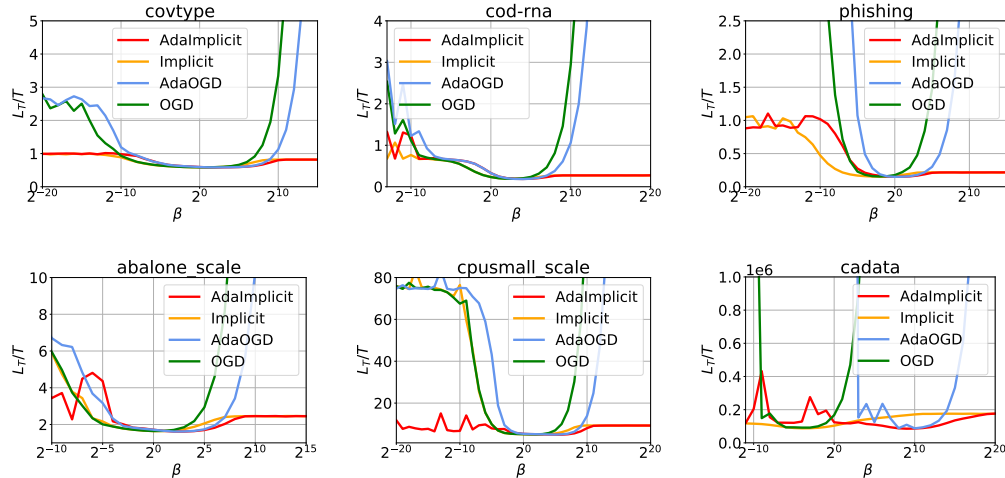


Figure 2: Plots on classification tasks using the hinge loss (top) and regression tasks using the absolute loss (bottom).

if AdaImplicit can achieve a good worst-case performance. We consider both classification and regression tasks. Additional plots can be found in Appendix D.

We used datasets from the LIBSVM library [6]. Before running the algorithms, we preprocess the data by dividing each feature by its maximum absolute value so that all the values are in the range $[-1, 1]$, then we add a bias term. Details about the datasets can be found in Appendix D.

Given that in the online setting we cannot tune the hyperparameter β using hold-out data, we plot the average cumulative loss of each algorithm, i.e., $L_t/t = \frac{1}{t} \sum_{i=1}^t \ell_i(\mathbf{x}_i)$, as a function of the hyperparameter β . This allows us to evaluate at the same time the sensitivity of the algorithms to β and their best performance with oracle tuning. Note that in all the algorithms we consider the optimal worst-case setting of β is proportional to the diameter of the feasible set, hence it is fair to plot their performance as a function of β . We consider values of β in $[2^{-20}, 2^{20}]$ with a grid containing 41 points. Then, each algorithm is run 10 times and results are averaged. For classification tasks we use the hinge loss, while for regression tasks we use the absolute loss. In both cases, we adopt the squared L_2 function for ψ . The details about implicit updates are discussed in Appendix C.

Results are illustrated in Fig. 2. From the plots, we can see that when fine-tuned, all the algorithms achieve similar results, i.e., the minimum value of average cumulative loss is very close for all the algorithms considered and there is not a clear winner. However, note that the range of values which allows an algorithm to reach the minimum is considerably wider for Implicit algorithms and confirms their robustness regarding learning rate misspecification, as already investigated in other works [see, e.g., 33, 34]. This is a great advantage when considering online algorithms since, contrarily to the batch setting, algorithms cannot be fine-tuned in advance relying on training/validation sets.

8 Conclusions

In this paper, we investigated online Implicit algorithms from a theoretical perspective. Our analysis revealed interesting insights regarding the behavior of these algorithms and allowed us to design a new *adaptive* algorithm, which may take advantage of “easy” data. The obtained experimental results indicate that in real-world tasks (such as online classification with hinge loss or online regression with the absolute loss), Implicit algorithms provide a better solution in terms of robustness, which is particularly relevant in online settings. Future directions include extending our analysis to a broader area, for example considering dynamic environments or strongly-convex loss functions, to see if the same gains can be proved. Finally, other examples of “easy” data can be considered, such as the case of stochastic loss functions.

Broader Impact

We believe our investigation will foster further studies promoting the adoption of adaptive learning rates in online learning and beyond. Indeed, in recent years adaptive methods in optimization proved to be one of the preferred methods for training deep neural networks. On the other hand, this work confirm the robustness of implicit updates and opens up to new possibilities in this field. From a societal aspect, this work is mainly theoretical and does not present any foreseeable consequence.

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Algorithm 3 OMD with Implicit updates

Require: Non-empty closed convex set $V \subset X \subset \mathbb{R}^d$, distance generating function $\psi : \mathbb{R}^d \rightarrow \mathbb{R}$, Lipschitz constant $L > 0$, $\beta > 0$, $\mathbf{x}_1 \in V$

- 1: Initialize: $i = 0$, $S_0 = 0$, $\eta_0 = \frac{\beta}{L}$
- 2: **for** $t = 1, \dots, T$ **do**
- 3: Output \mathbf{x}_t
- 4: Receive $\ell_t : \mathbb{R}^d \rightarrow (-\infty, +\infty]$ and pay $\ell_t(\mathbf{x}_t)$
- 5: $\mathbf{x}'_{t+1} = \arg \min_{\mathbf{x} \in V} B_\psi(\mathbf{x}, \mathbf{x}_t) + \eta_i \ell_t(\mathbf{x})$
- 6: $\delta_t = \ell_t(\mathbf{x}_t) - \ell_t(\mathbf{x}'_{t+1}) - \frac{B_\psi(\mathbf{x}'_{t+1} - \mathbf{x}_t)}{\eta_i}$
- 7: $S_i \leftarrow S_i + \delta_t$
- 8: **if** $S_i \geq \eta_i L^2 2^i$ **then**
- 9: $i = i + 1$
- 10: Update $\eta_i = \frac{\beta}{L\sqrt{2^i}}$
- 11: $S_i = 0$
- 12: $\mathbf{x}_{t+1} = \mathbf{x}_1$
- 13: **else**
- 14: $\mathbf{x}_{t+1} = \mathbf{x}'_{t+1}$
- 15: **end if**
- 16: **end for**

A Doubling Trick

In this section, we present a doubling trick strategy to tune the learning rate in IOMD. As already mentioned in Section 6, it is possible to apply this construction if we consider the terms δ_t , whose sum over time is an increasing sequence thanks to Eq. (5). We present the doubling trick here to show that it is not simpler than our main analysis in the paper nor it adds empirical advantages. The only advantage coming from using a doubling trick derives from the fact that it is not required to have a bounded domain (as we will show in Theorem A.2), opposed to the case of an adaptive learning rate.

In the following, we slightly modify the definition of δ_t given in Eq. (14) (see line 6 of Algorithm 3),

$$\delta_t = \ell_t(\mathbf{x}_t) - \ell_t(\mathbf{x}'_{t+1}) - \frac{B_\psi(\mathbf{x}_t, \mathbf{x}'_{t+1})}{\eta_t}.$$

The algorithm works as follows: at the beginning of epoch i , we set the learning rate $\eta_i = \beta/L\sqrt{2^i}$, run IOMD and monitor the sum $\sum_t \delta_t/\eta_i$ until it reaches $L^2 2^i$. Once this happens, we restart the algorithm doubling the threshold and halving the learning rate—see Algorithm 3. Note that during an epoch the learning rate stays fixed. Let t_i be the index of the first round of epoch i . We then have $\eta_t = \eta_i$ for $t \in [t_i, t_{i+1} - 1]$.

For the analysis, note that by using a doubling trick the time horizon is divided in N different epochs (where N is obviously not known a priori). We next show that the total number of epochs N is logarithmic in a quantity Δ_i which we define for the sake of the analysis.

Lemma A.1. *Let t_i be the first time-step of epoch i , with $t_0 = 1$. Suppose Algorithm 3 is run for a total of N epochs. Let $\Delta_i \triangleq \sum_{t=t_i}^{t_{i+1}-1} \delta_t$. Then, we have that*

$$N \leq \min \left(\log_2 \left(\frac{1}{L^2} \sum_{i=0}^N \frac{\Delta_i}{\eta_i} + 1 \right), 2 \log_2 \left(\frac{\sqrt{2}-1}{L} \sum_{i=0}^N \Delta_i + 1 \right) \right). \quad (17)$$

Proof. First, we have that

$$\sum_{i=0}^{N-1} a^i = \frac{a^N - 1}{a - 1}. \quad (18)$$

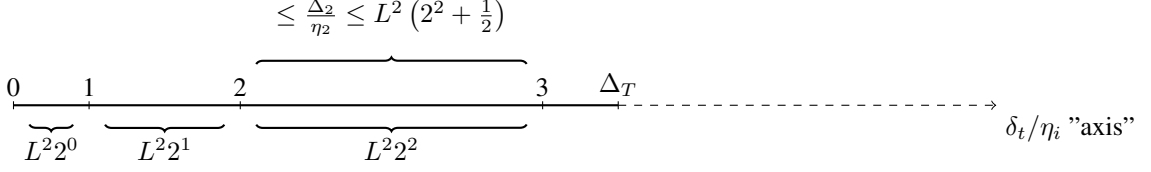


Figure 3: Doubling trick illustrated.

Now, note that the sum of the $L^2 2^i$ terms in the first $N - 1$ epochs is less than or equal to the final sum of the terms monitored. Therefore, using Eq. (18) we have

$$\sum_{i=0}^{N-1} L^2 2^i = L^2 (2^N - 1) \leq \sum_{i=0}^N \sum_{t=t_i}^{t_{i+1}-1} \frac{\delta_t}{\eta_i} = \sum_{i=0}^N \frac{\Delta_i}{\eta_i},$$

and solving for N yields the first term in the min in Eq. (17).

On the other hand, we have that the condition for the restart can be rewritten as $\sum_t \delta_t \leq L\sqrt{2^i}$. Using again Eq. (18), we get

$$\sum_{i=0}^{N-1} L\sqrt{2^i} = L \frac{2^{\frac{N}{2}} - 1}{\sqrt{2} - 1} \leq \sum_{i=0}^N \sum_{t=t_i}^{t_{i+1}-1} \delta_t = \sum_{i=0}^N \Delta_i, \quad (19)$$

rearranging and solving for N gives the second term in the min in Eq. (17). \square

We can now analyze the regret incurred by the algorithm. Each epoch is bounded individually. The final regret bound is given by the sum of the individual contributions over all the epochs.

Theorem A.2. *Assume the losses $\ell_t(\mathbf{x})$ to be L -Lipschitz, for all $t = 1, \dots, T$ and ψ to be 1-strongly convex w.r.t. $\|\cdot\|$. Let Δ_i be defined as in Lemma A.1. Then, for any $\mathbf{u} \in \mathbb{R}^d$ the regret of Algorithm 3 after T rounds is bounded as*

$$R_T(\mathbf{u}) \leq c \left(\frac{B_\psi(\mathbf{u}, \mathbf{x}_1)}{\beta} + \beta \right) L\sqrt{T+1} + c \frac{\beta L}{2}, \quad (20)$$

where $c = \frac{\sqrt{2}}{\sqrt{2}-1}$.

Proof. Using the notation introduced, we have that $\delta_{t_{i+1}-1}$ is the term which causes the restart in epoch i . Therefore, the following holds

$$\frac{\Delta_i}{\eta_i} \leq L^2 \left(2^i + \frac{1}{2} \right), \quad (21)$$

since from Lemma 5.3 the last term in epoch i is such that $\frac{\delta_{t_{i+1}-1}}{\eta_i} \leq \frac{L^2}{2}$.

We next define $R_i(\mathbf{u})$ as the regret during epoch i ,

$$R_i(\mathbf{u}) = \sum_{t=t_i}^{t_{i+1}-1} (\ell_t(\mathbf{x}_t) - \ell_t(\mathbf{u})).$$

Using Eq. (10) with $\eta_t = \eta_i$ for $t \in [t_i, \dots, t_{i+1} - 1]$, we have that $R_i(\mathbf{u})$ is bounded as follows

$$\begin{aligned} R_i(\mathbf{u}) &\leq \frac{B_\psi(\mathbf{u}, \mathbf{x}_{t_i})}{\eta_i} + \sum_{t=t_i}^{t_{i+1}-1} \left[\ell_t(\mathbf{x}_t) - \ell_t(\mathbf{x}'_{t+1}) - \frac{1}{\eta_i} B_\psi(\mathbf{x}'_{t+1}, \mathbf{x}_t) \right] \\ &\leq \frac{B_\psi(\mathbf{u}, \mathbf{x}_{t_i})}{\eta_i} + \sum_{t=t_i}^{t_{i+1}-1} \delta_t = \frac{B_\psi(\mathbf{u}, \mathbf{x}_1)}{\eta_i} + \Delta_i. \end{aligned} \quad (22)$$

We can now write the final regret bound by summing $R_i(\mathbf{u})$ over all the epochs $i = 1, \dots, N$.

$$\begin{aligned}
R_T(\mathbf{u}) &= \sum_{i=0}^N R_i(\mathbf{u}) \\
&\leq \sum_{i=0}^N \left[\frac{B_\psi(\mathbf{u}, \mathbf{x}_1)}{\eta_i} + \Delta_i \right] \\
&\leq \sum_{i=0}^N \left[\frac{B_\psi(\mathbf{u}, \mathbf{x}_1)}{\eta_i} + \eta_i L^2 \left(2^i + \frac{1}{2} \right) \right] && \text{see Eq. (21)} \\
&= L \left(\frac{B_\psi(\mathbf{u}, \mathbf{x}_1)}{\beta} + \beta \right) \sum_{i=0}^N \sqrt{2^i} + \frac{\beta L}{2} \sum_{i=0}^N \frac{1}{\sqrt{2^i}} && \left(\eta_i = \frac{\beta}{L\sqrt{2^i}} \right) \\
&\leq \left(\frac{B_\psi(\mathbf{u}, \mathbf{x}_1)}{\beta} + \beta \right) L \frac{2^{\frac{N+1}{2}} - 1}{\sqrt{2} - 1} + \frac{\beta L}{2} \frac{\sqrt{2}}{\sqrt{2} - 1} \\
&\leq c \left(\frac{B_\psi(\mathbf{u}, \mathbf{x}_1)}{\beta} + \beta \right) L 2^{\frac{N}{2}} + c \frac{\beta L}{2}, && (23)
\end{aligned}$$

where in the last step we used Eq. (18) and the definition of c .

Now, using the first term in Eq. (17) we have that

$$2^{\frac{N}{2}} = 2^{\frac{1}{2} \log_2 \left(\frac{1}{L^2} \sum_{i=0}^N \frac{\Delta_i}{\eta_i} + 1 \right)} = \sqrt{\frac{1}{L^2} \sum_{i=0}^N \frac{\Delta_i}{\eta_i} + 1}.$$

Therefore, from Eq. (23) the regret can be bounded as

$$R_T(\mathbf{u}) \leq c \left(\frac{B_\psi(\mathbf{u}, \mathbf{x}_1)}{\beta} + \beta \right) \sqrt{\sum_{i=0}^N \frac{\Delta_i}{\eta_i} + L^2} + c \frac{\beta L}{2}.$$

Furthermore, using Theorem 5.3 and the assumption on the losses to be L -Lipschitz, we get

$$\sum_{i=0}^N \frac{\Delta_i}{\eta_i} = \sum_{i=0}^N \sum_{t=t_i}^{t_{i+1}-1} \frac{\delta_t}{\eta_i} \leq \sum_{i=0}^N \sum_{t=t_i}^{t_{i+1}-1} \frac{L^2}{2} = \frac{L^2}{2} T,$$

substituting back the above result we get

$$R_T(\mathbf{u}) \leq cL \left(\frac{B_\psi(\mathbf{u}, \mathbf{x}_1)}{\beta} + \beta \right) \sqrt{T+1} + c \frac{\beta L}{2}. \quad \square$$

In principle, it is possible to get a bound which interpolates between $\mathcal{O}(\sqrt{T})$ and $\mathcal{O}(V_T)$ as done in Theorem 6.2. However, for the latter possibility a bounded domain seems required, in order to bound the difference between the first and last losses of each epoch. Moreover, the number of restarts is in the worst case order of $\mathcal{O}(\ln T)$, which would give a (slightly) worse bound compared to Theorem 6.2. For these reasons, we will not provide details about how to use Algorithm 3 in order to get a regret bound order of $\mathcal{O}(V_T)$. Nonetheless, we next show how in the case of fixed losses, we can still recover a constant regret bound, even if the domain is unbounded.

A.1 Fixed Losses

If the losses are all equal, i.e. $\ell_t(\mathbf{x}) = \ell(\mathbf{x})$ for all $t = 1, \dots, T$, running Algorithm 3 from Theorem A.2 we would expect a regret which scales as $\mathcal{O}(\sqrt{T})$. However, as we are going to show in the next lemma, with a proper setting of β this is actually not the case and Algorithm 3 always stays in the first epoch, even if the domain is unbounded!

Lemma A.3. *Assume the losses over time are fixed, i.e. $\ell_t(\mathbf{x}) = \ell(\mathbf{x})$ for all $t = 1, \dots, T$ and $\psi(\mathbf{x})$ be 1-strongly convex w.r.t $\|\cdot\|$. Then, Algorithm 3 with $\beta \geq 1$ will always stay in the first epoch, i.e., $N = 0$.*

Proof. In order to not have a restart, we need $\delta_t \leq \eta_i L^2 2^i$, which translates to

$$\ell(\mathbf{x}_t) - \ell(\mathbf{x}'_{t+1}) \leq \beta L \sqrt{2^i} (B_\psi(\mathbf{x}'_{t+1}, \mathbf{x}_t) + 1) .$$

From the inequality above, we have a reset iff

$$\begin{aligned} i &\leq 2 \log_2 \frac{\ell(\mathbf{x}_t) - \ell(\mathbf{x}'_{t+1})}{\beta L (B_\psi(\mathbf{x}'_{t+1}, \mathbf{x}_t) + 1)} \leq 2 \log_2 \frac{L \|\mathbf{x}_t - \mathbf{x}'_{t+1}\|}{\beta L (1 + \frac{1}{2} \|\mathbf{x}'_{t+1} - \mathbf{x}_t\|^2)} \\ &= 2 \log_2 \frac{2 \|\mathbf{x}'_{t+1} - \mathbf{x}_t\|}{\beta (2 + \|\mathbf{x}'_{t+1} - \mathbf{x}_t\|^2)}, \end{aligned}$$

where the second inequality derives from the Lipschitzness of the losses and by lower bounding the Bregman divergence with Eq. (2). Now, let $f(y) = \frac{2y}{2+y^2}$, with $y \geq 0$. We have that $\lim_{y \rightarrow +\infty} f(y) = 0$ and $f(0) = 0$. Furthermore, if we take the derivative and set it to 0, we see that $f(y)$ has a maximum in $y = \sqrt{2}$. Hence, we have

$$i \leq 2 \log_2 \frac{2 \|\mathbf{x}'_{t+1} - \mathbf{x}_t\|}{\beta (2 + \|\mathbf{x}'_{t+1} - \mathbf{x}_t\|^2)} \leq 2 \log_2 \frac{1}{\beta \sqrt{2}} = -1 - 2 \log_2 \beta \leq -1,$$

where the last step derives from the assumption on β . Therefore, if $i > -1$ then we will not double the learning rate. Note that this is always verified and we can conclude that $N = 0$. \square

We can now prove a regret bound in the case of fixed losses.

Theorem A.4. *Under the assumptions of lemma A.3 the regret incurred by Algorithm 3 is bounded as*

$$R_T(\mathbf{u}) \leq \frac{L}{\beta} B_\psi(\mathbf{u}, \mathbf{x}_1) + \ell(\mathbf{x}_1) - \ell(\mathbf{x}_T) . \quad (24)$$

Proof. From Eq. (22) and the fact that we only have one epoch thanks to lemma A.3, we have that

$$\begin{aligned} R_T(\mathbf{u}) = R_0(\mathbf{u}) &\leq \frac{B_\psi(\mathbf{u}, \mathbf{x}_1)}{\eta_0} + \sum_{t=1}^T \delta_t \\ &= \frac{B_\psi(\mathbf{u}, \mathbf{x}_1)}{\eta_0} + \sum_{t=1}^T \left[\ell(\mathbf{x}_t) - \ell(\mathbf{x}_{t+1}) - \frac{B_\psi(\mathbf{x}_{t+1}, \mathbf{x}_t)}{\eta_0} \right] \\ &\leq \frac{B_\psi(\mathbf{u}, \mathbf{x}_1)}{\eta_0} + \ell(\mathbf{x}_1) - \ell(\mathbf{x}_T) \\ &= \frac{L}{\beta} B_\psi(\mathbf{u}, \mathbf{x}_1) + \ell(\mathbf{x}_1) - \ell(\mathbf{x}_T) . \quad \square \end{aligned}$$

B Proofs

The proof of the properties in Proposition 4.1 is straightforward, but we report it here for completeness.

Proof of Proposition 4.1. From the update in Eq. (4), we immediately get the following inequality

$$\eta_t \ell_t(\mathbf{x}_{t+1}) \leq \eta_t \ell_t(\mathbf{x}_{t+1}) + B_\psi(\mathbf{x}_{t+1}, \mathbf{x}_t) \leq \eta_t \ell_t(\mathbf{x}_t) + B_\psi(\mathbf{x}_t, \mathbf{x}_t) = \eta_t \ell_t(\mathbf{x}_t),$$

which verifies Eq. (5) and implies that value of the loss ℓ_t in \mathbf{x}_{t+1} is not bigger than the one in \mathbf{x}_t .

Eq. (6) is simply the first-order optimality condition for \mathbf{x}_{t+1} .

For Eq. (7), from the convexity of ℓ_t we have

$$\begin{aligned} \ell_t(\mathbf{x}_{t+1}) &\geq \ell_t(\mathbf{x}_t) + \langle \mathbf{g}_t, \mathbf{x}_{t+1} - \mathbf{x}_t \rangle, \\ \ell_t(\mathbf{x}_t) &\geq \ell_t(\mathbf{x}_{t+1}) + \langle \mathbf{g}'_t, \mathbf{x}_t - \mathbf{x}_{t+1} \rangle . \end{aligned}$$

Summing both inequalities, we get the desired result. \square

Proof of Lemma 5.1. From Eq. (6), we have that

$$\begin{aligned}\eta_t(\ell_t(\mathbf{x}_{t+1}) - \ell_t(\mathbf{u})) &\leq \langle \eta_t \mathbf{g}'_t, \mathbf{x}_{t+1} - \mathbf{u} \rangle \leq \langle \nabla \psi(\mathbf{x}_t) - \nabla \psi(\mathbf{x}_{t+1}), \mathbf{x}_{t+1} - \mathbf{u} \rangle \\ &= B_\psi(\mathbf{u}, \mathbf{x}_t) - B_\psi(\mathbf{u}, \mathbf{x}_{t+1}) - B_\psi(\mathbf{x}_{t+1}, \mathbf{x}_t).\end{aligned}$$

Dividing by η_t and summing over $t = 1, \dots, T$ yields the result in Eq. (8).

For the second part observe that

$$\sum_{t=1}^T \frac{B_\psi(\mathbf{u}, \mathbf{x}_t) - B_\psi(\mathbf{u}, \mathbf{x}_{t+1})}{\eta_t} \leq \frac{D^2}{\eta_1} + D^2 \sum_{t=2}^T \left(\frac{1}{\eta_t} - \frac{1}{\eta_{t-1}} \right) = \frac{D^2}{\eta_T}. \quad \square$$

Proof of Lemma 6.1. From the assumptions, we have that $\Delta_{t+1} \leq \Delta_t + \min\{ba_t, ca_t^2/(2\Delta_t)\}$, for any $t \geq 1$. Also, observe that

$$\Delta_{T+1}^2 = \sum_{t=1}^T \Delta_{t+1}^2 - \Delta_1^2 = \sum_{t=1}^T \left[\underbrace{(\Delta_{t+1} - \Delta_t)^2}_{(a)} + 2 \underbrace{(\Delta_{t+1} - \Delta_t)\Delta_t}_{(b)} \right].$$

We bound (a) and (b) separately. For (a), from the assumption on the recurrence and using the first term in the minimum we have that $(\Delta_{t+1} - \Delta_t)^2 \leq b^2 a_t^2$. On the other hand, for (b) using the second term in the minimum in the recurrence we get $2(\Delta_{t+1} - \Delta_t)\Delta_t \leq ca_t^2$. Putting together the results we have that $\Delta_{T+1}^2 \leq (b^2 + c) \sum_{t=1}^T a_t^2$ and the lemma follows. \square

B.1 Lower Bound

Proof of Theorem 6.3. The first loss of the algorithm is $\ell_1(\mathbf{x}) = L\langle \mathbf{g}, \mathbf{x} \rangle$, where $\|\mathbf{g}\|_* = 1$ and \mathbf{g} is orthogonal to \mathbf{x}_1 , while L will be set in the following. Note that $d \geq 2$ assures that \mathbf{g} always exists. For $t \geq 2$, set $\ell_t(\mathbf{x}) = 0$. First, observe that $V_T = \max_{\mathbf{x} \in V} -\ell_1(\mathbf{x}) = L \max_{\mathbf{u} \in V} -\langle \mathbf{g}, \mathbf{u} \rangle = LD/2$. Hence, setting $L = 2V'_T/D$, we have $V_T = V'_T$. Also, we have $R_T(\mathbf{u}) = L \max_{\mathbf{u} \in V} -\langle \mathbf{g}, \mathbf{u} \rangle = V_T$. \square

It is worth emphasizing that the lower bound does not contradict the upper bound of $\mathcal{O}(DL\sqrt{T})$ because here L is chosen arbitrarily large.

C Formulas

First, let's mention the update rules for IOMD with $V = \mathbb{R}^d$ for hinge loss, absolute loss, and square loss respectively [see, e.g., 11, 18].

$$\begin{aligned}\mathbf{x}_{t+1} &= \mathbf{x}_t + \min \left(\eta_t, \frac{\max(1 - y_t \langle \mathbf{z}_t, \mathbf{x}_t \rangle, 0)}{\|\mathbf{z}_t\|^2} \right) y_t \mathbf{z}_t, \\ \mathbf{x}_{t+1} &= \mathbf{x}_t - \min \left(\eta_t, \frac{|\langle \mathbf{z}_t, \mathbf{x}_t \rangle - y_t|}{\|\mathbf{z}_t\|^2} \right) \mathbf{z}_t, \\ \mathbf{x}_{t+1} &= \mathbf{x}_t - \eta_t \frac{(\langle \mathbf{z}_t, \mathbf{x}_t \rangle - y_t) \mathbf{z}_t}{1 + \eta \|\mathbf{z}_t\|_2^2}.\end{aligned}$$

Now let's consider the case that $V = \{\mathbf{x} : \|\mathbf{x}\|_2 \leq D/2\}$. In this case it is easy to see that the update becomes

$$\mathbf{x}_{t+1} = \mathbf{x}_t - \eta_t \mathbf{g}'_t - \alpha \mathbf{x}_{t+1},$$

where $\mathbf{g}'_t \in \partial \ell_t(\mathbf{x}_{t+1})$, $\alpha \mathbf{x}_{t+1} \in \partial i_V(\mathbf{x}_{t+1})$ and $\alpha \geq 0$. Hence, we have that

$$\mathbf{x}_{t+1} = \frac{\mathbf{x}_t}{\alpha + 1} - \frac{\eta_t}{\alpha + 1} \mathbf{g}'_t.$$

This implies that we can take the previous formulas and substitute $\frac{\mathbf{x}_t}{\alpha + 1}$ to \mathbf{x}_t and $\frac{\eta_t}{\alpha + 1}$ to η_t . Then, the optimal $\alpha \geq 0$ is the smallest one that gives $\|\mathbf{x}_{t+1}\| \leq D/2$. Note that this is a 1-dimensional problem that can be easily solved numerically.

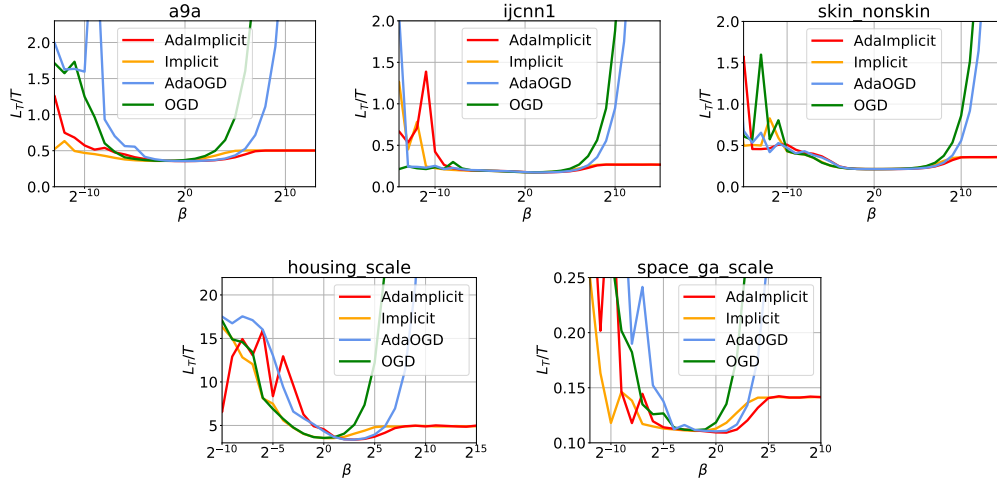


Figure 4: Plots on classification tasks using the hinge loss (top) and regression tasks using the absolute loss (bottom).

D Experiments

In Fig. 4 we show plots about other experiments on real data which were not shown in the main paper. Details about the datasets used can be found in Table 1 and Table 2.

Table 1: Classification datasets

Name	Datapoints	Features
a9a	32,561	123
ijcnn1	49,990	22
cod-rna	59,535	8
covtype	581,012	54
skin_nonskin	245,057	3
phishing	11,055	68

Table 2: Regression datasets

Name	Datapoints	Features
abalone	11,055	8
cadata	20,640	8
cpusmall	8,192	12
housing	506	13
space_ga	3,107	6

E Implicit Updates for FTRL

In this section, we show how to get a bound of $\mathcal{O}(V_T + 1)$ for FTRL employed with full losses. As already explained in the main paper, contrarily to the OMD case, we do not have efficient algorithms to solve the minimization problem given by the FTRL update rule in this case. Furthermore, we show that it is not possible to adopt the same learning rate tuning strategy of *Adalmplicit* in order to get a similar regret bound.

We first remember the FTRL regret bound, which is standard in the literature (see e.g. [24] Lemma 7.1).

Theorem E.1. *Let $V \subseteq \mathbb{R}^d$ be closed and non-empty. Denote by $F_t(\mathbf{x}) = \psi_t(\mathbf{x}) + \sum_{i=1}^{t-1} \ell_i(\mathbf{x})$, where ψ_1, \dots, ψ_T is a sequence of regularizers such that $\psi_t : \mathbb{R}^d \rightarrow (-\infty, +\infty]$ for all t . Assume that $\operatorname{argmin}_{\mathbf{x} \in V}$ is not empty and set $\mathbf{x}_1 = \operatorname{argmin}_{\mathbf{x} \in V} \psi_1(\mathbf{x})$, $\mathbf{x}_{t+1} \in \operatorname{argmin}_{\mathbf{x} \in V} F_t(\mathbf{x})$. Then, for any \mathbf{u} , we have*

$$R_T(\mathbf{u}) \leq (\psi_T(\mathbf{u}) - \psi_T(\mathbf{x}_1)) + \sum_{t=1}^T [F_t(\mathbf{x}_t) - F_{t+1}(\mathbf{x}_{t+1}) + \ell_t(\mathbf{x}_t)]. \quad (25)$$

Now, assume that $\psi_t(\mathbf{x}) = \lambda_t \psi(\mathbf{x})$, with $(\lambda_t)_{t=1}^T$ being a non-decreasing sequence. We can rewrite the sum over time on the right-hand side of Eq. (25) as follows

$$\begin{aligned}
& \sum_{t=1}^T [F_t(\mathbf{x}_t) - F_{t+1}(\mathbf{x}_{t+1}) + \ell_t(\mathbf{x}_t)] \\
&= \sum_{t=1}^T [F_t(\mathbf{x}_t) + \ell_t(\mathbf{x}_t) - F_t(\mathbf{x}_{t+1}) - \ell_t(\mathbf{x}_{t+1}) + \psi_t(\mathbf{x}_{t+1}) - \psi_{t+1}(\mathbf{x}_{t+1})] \\
&\leq \sum_{t=1}^T [\ell_t(\mathbf{x}_t) - \ell_t(\mathbf{x}_{t+1}) + \psi_t(\mathbf{x}_{t+1}) - \psi_{t+1}(\mathbf{x}_{t+1})] \\
&= \ell_1(\mathbf{x}_1) - \ell_T(\mathbf{x}_{T+1}) + \sum_{t=2}^T [\ell_t(\mathbf{x}_t) - \ell_{t-1}(\mathbf{x}_t)] + \sum_{t=1}^T (\lambda_t - \lambda_{t+1}) \psi(\mathbf{x}_{t+1}) \\
&\leq \ell_1(\mathbf{x}_1) - \ell_T(\mathbf{x}_{T+1}) + V_T + \sum_{t=1}^T (\lambda_t - \lambda_{t+1}) \psi(\mathbf{x}_{t+1}),
\end{aligned}$$

where the first inequality derives from the fact that $F_t(\mathbf{x}_t) - F_t(\mathbf{x}_{t+1}) \leq 0$ while the last one from the definition of V_T . Therefore, the regret bound can be rewritten as follows

$$R_T(\mathbf{u}) \leq \lambda_T(\psi(\mathbf{u}) - \psi(\mathbf{x}_1)) + \ell_1(\mathbf{x}_1) - \ell_T(\mathbf{x}_{T+1}) + V_T + \sum_{t=1}^T (\lambda_t - \lambda_{t+1}) \psi(\mathbf{x}_{t+1})$$

We can see that with a constant learning rate the above expression would give a regret bound $\mathcal{O}(V_T + 1)$. On the other hand, it is known from the literature that a parameter $\lambda_t \propto \sqrt{t}$ would give a regret bound of $\mathcal{O}(\sqrt{T})$. Ideally, we would like to have a certain λ_t which allows to interpolate between a regret bound of $\mathcal{O}(V_T + 1)$ and $\mathcal{O}(\sqrt{T})$, as done for *AdaImplicit*. However, the techniques adopted in Section 6 do not seem to work in this case and one should hence resort to a different approach. In addition to the technical difficulties, as already stated in the main paper the computational burden of implicit updates with FTRL could be prohibitive in practice and makes this approach not worth of pursuing.