

1 We thank the reviewers for their detailed and constructive comments. Please find our answers below.

Algorithm	Lazy TS	RAGE	Oracle	Algorithm	Lazy TS	RAGE	$\mathcal{XY}$ -Adaptive	Oracle
dimension	Mean (Std)	Mean (Std)	Mean (Std)	confidence level	Mean (Std)	Mean (Std)	Mean (Std)	Mean (Std)
$(d = 4)$	<b>3346.29</b> (125.3)	8033.51 (464.0)	3968.72 (163.9)	$(\delta = 0.5)$	<b>3080.56</b> (119.1)	5840.21 (373.6)	6192.34 (373.8)	3016.91 (133.1)
$(d = 7)$	<b>4405.74</b> (143.5)	9675.00 (537.8)	4107.09 (160.6)	$(\delta = 0.05)$	<b>3699.84</b> (130.1)	7751.79 (434.2)	8167.51 (368.3)	3610.33 (146.1)
$(d = 10)$	<b>5602.55</b> (180.9)	9780.54 (360.0)	4321.84 (167.8)	$(\delta = 0.005)$	<b>4297.23</b> (131.8)	9810.19 (543.8)	9278.82 (315.8)	4219.38 (165.3)

2 **Reviewer 1. 1) Experiments.** An important part of our work is the design of the first optimal algorithm whose  
3 implementation and performance guarantees are completely independent of the number  $K$  of arms. That is why our  
4 experiments focused on scenarios with large set of arms. We will include more experiments (essentially all scenarios  
5 considered in [16]) – e.g. the table above corresponds to the benchmark used in Soare et al. [12] and all other papers  
6 on the topic: here the angle  $\omega = 0.1$ , the left part of the table is for  $\delta = 0.01$ , and the right part for  $d = 6$ . In all  
7 experiments we have done, the results suggest that our algorithm is very competitive.

8 **2) Explore-and-Verify (EV) framework.** Thanks for suggesting that combining RAGE and the EV framework of  
9 [13] may lead to an efficient algorithm. As you mention, an idea of the same flavour is used by Chen et al. [0] for  
10 combinatorial pure exploration. But to get guarantees in expectation, one needs parallel simulations (Lemma 4.8  
11 in [0]) or repetitive calls of RAGE (see Appendix 2 of [13]). Thus, as already noticed by Fiez et al., the authors of  
12 RAGE, the EV framework may be impractical (see [16] Page 7). In fact, neither the authors of [0] nor those of [13]  
13 implemented their framework in practice. In our paper, we devise a very simple and asymptotically optimal algorithm  
14 using a track-and-stop framework in the spirit of [18]. To this aim, we needed to tackle the following challenges: (i) the  
15 optimal allocation is not unique, which poses tracking issues; (ii) computing an optimal allocation in each round might  
16 be computationally demanding (we solve this issue with the laziness of our tracking rule); (iii) most of the components  
17 of the algorithm, including the stopping rule, must be independent of the number of arms  $K$ .

18 **3) Continuous sets of arms.** Our main contribution for such sets of arms is to derive the first (problem-specific) sample  
19 complexity lower bound at least for the sphere. This has not been done earlier, and proved to be very challenging. It  
20 might be the case that discretizing the set of arms using an  $\epsilon/2$ -net of the sphere and applying our TS algorithm would  
21 yield a sample complexity with the right scaling in  $d, \epsilon, \delta$ . The algorithm we propose is even simpler, and in particular  
22 does not need to work with a set of arms with cardinality growing exponentially with the dimension.

23 **Reviewer 3. 1) Novelty of the problem and of our solution.** The problem of best arm identification in linear bandits  
24 is not new, but it is fundamental. To the best of our knowledge, we propose the first asymptotically optimal algorithm  
25 whose implementation and performance guarantees are completely independent of the number  $K$  of arms. To this aim  
26 and to use the track-and-stop framework [18], we needed to tackle the following challenges: (i) the optimal allocation is  
27 not unique, which poses tracking issues; (ii) computing an optimal allocation in each round might be computationally  
28 demanding (we solve this issue with the laziness of our tracking rule); (iii) most of the components of the algorithm,  
29 including the stopping rule, must be independent of the number of arms  $K$ . We agree that a very interesting research  
30 direction is to derive non-asymptotic performance guarantees for our algorithm.

31 **Reviewer 4.** Thanks for pointing us to the recent paper of Degenne et al. (it was not available when we submitted our  
32 paper). We will cite and discuss it. We also appreciate your suggestions for improving the clarity of the paper. We will  
33 elaborate further on the broader impact section as suggested.

34 **1) Experiments.** Please refer to answer 1) to Reviewer 1.

35 **2) The constant  $c_{\mathcal{A}_0}$ .** This constant controls the rate at which forced exploration is performed. Having a too large  
36 constant could indeed lead to worse performance. Theoretically, however, it does not affect the asymptotic performance  
37 of the algorithm. In practice, we note that forced exploration rarely occurs for the considered level of laziness in our  
38 experiments. We conjecture that it may not even be needed. We will provide further discussion on this in the paper.

39 **3) The goal of  $\mathcal{A}_0$ .** We confirm that the goal of  $\mathcal{A}_0$  is to ensure that the growth rate of  $\lambda_{\min}(\sum_{s=1}^t a_s a_s^\top)$  is not too  
40 small. We also note that the chosen growth rate  $f(t) = O(\sqrt{t})$  is not too large because, ideally, if we knew an optimal  
41 allocation, sampling according to it would yield a growth rate of order exactly  $t$ .

42 **4) Estimation of  $\mu$  along problem-dependent directions.** These directions also appear implicitly in the sampling rule.  
43 More precisely, the tracked weights  $w_t$  in the sampling rule are solutions of the optimization problem  $\max_{w \in \Lambda} \psi(\hat{\mu}_t, w)$ .

44 **5) The stopping rule interpretation.** Intuitively, the stopping rule can be viewed as an empirical version of the problem-  
45 specific sample complexity lower bound. Roughly speaking, the stopping rule would correspond to the lower bound  
46 where  $\mu$  is replaced by  $\hat{\mu}_t$ , and  $w$  by  $(N_a(t)/t)_{a \in \mathcal{A}}$ .

47 **6) Continuous set of arms.** Regarding the continuous set of arms, our main result is the lower bound. Deriving such  
48 an explicit bound was challenging: in the change-of-measure argument, we needed to propose an appropriate reduction  
49 of the set of *confusing* parameters, see Appendix G, Steps 2 and 3 of the proof. We are only able to derive it for the  
50 sphere for now. For a generic continuous set of arms, the problem is even harder because in the change-of-measure  
51 argument, the aforementioned reduction of the set of confusing parameters becomes even more challenging.